

# Extensive-Form Games with Imperfect Information

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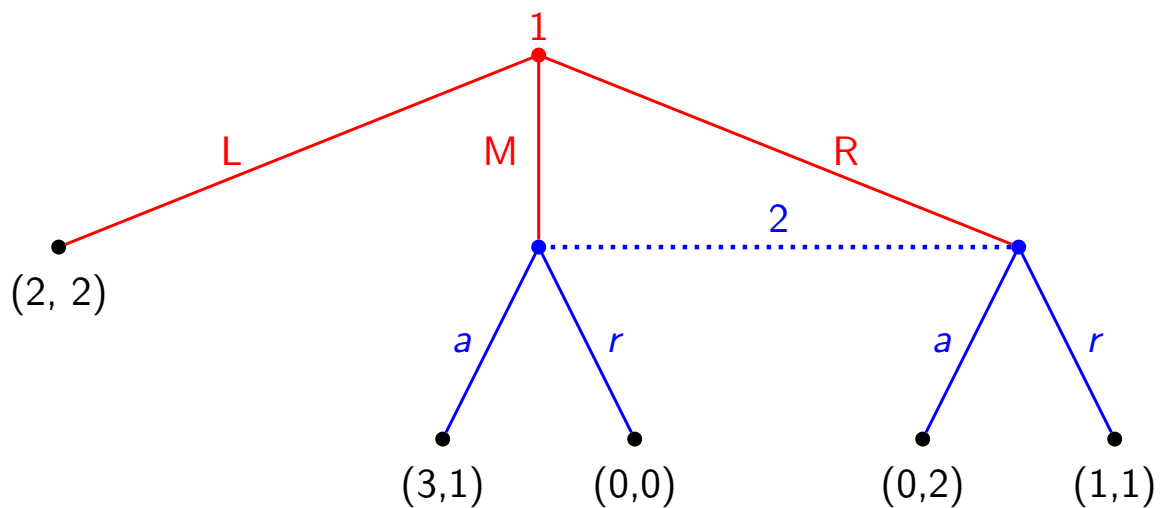
## Perfect Information vs. Imperfect Information

- ▶ Perfect Information
  - ▶ All players know the game structure.
  - ▶ Each player, when making any decision, **is perfectly informed** of all the events that have **previously** occurred.
- ▶ Imperfect Information
  - ▶ All players know the game structure.
  - ▶ Each player, when making any decision, may **not** be perfectly informed about some (or all) of the events that have already occurred.

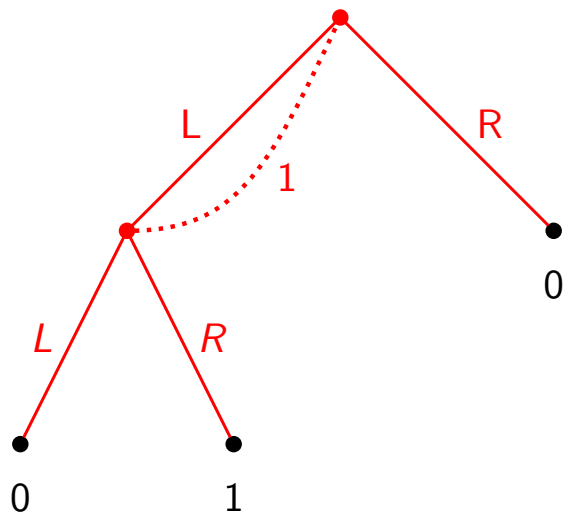
# Roadmap

- ▶ Define Imperfect-Information Extensive-Form Game
- ▶ Introduce **Sequential Equilibrium**  
*"rather a lot of bodies are buried in this definition". (Kreps, 1990)*

## Example 1



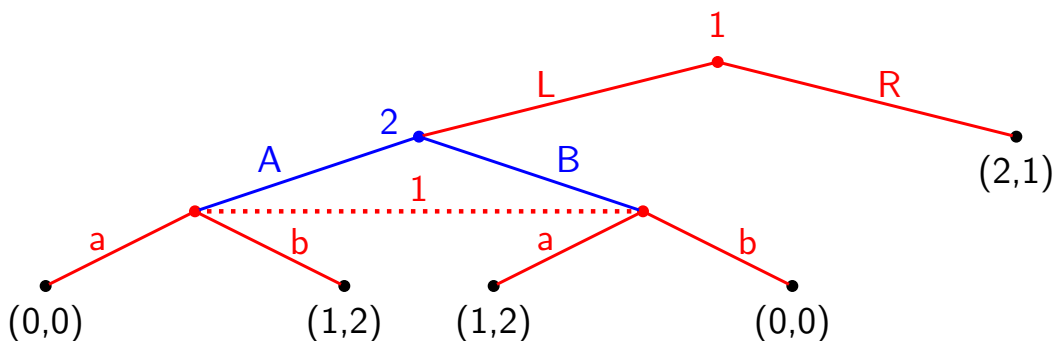
## Example 2



### Def. of Imperfect-Information Extensive-Form Games

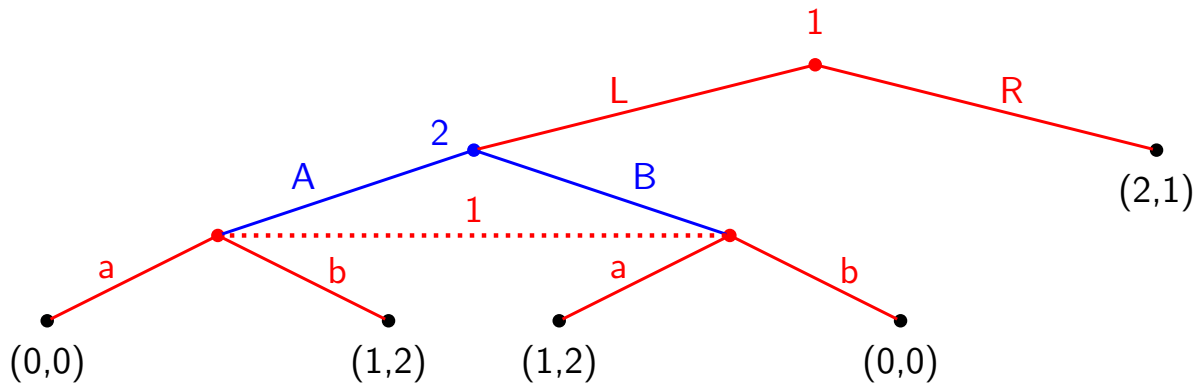
- ▶ An imperfect-information extensive-form game is a tuple  $(N, H, P, \mathcal{I}, u)$ 
  - ▶  $(N, H, P, u)$  is a perfect-information extensive-form game
  - ▶  $\mathcal{I} = \{\mathcal{I}_1, \mathcal{I}_2, \dots, \mathcal{I}_n\}$  is the set of information partitions of all players
    - ▶  $\mathcal{I}_i = \{\mathcal{I}_{i,1}, \dots, \mathcal{I}_{i,k_i}\}$  is the information partition of player  $i$
    - ▶  $\mathcal{I}_{i,j}$  is an information set of player  $i$
    - ▶ Action set  $A(h) = A(h')$  if  $h$  and  $h'$  are in  $\mathcal{I}_{i,j}$ , denote as  $A(\mathcal{I}_{i,j})$
    - ▶  $P(\mathcal{I}_{i,j})$  be the player who plays at information set  $\mathcal{I}_{i,j}$ .

$$\mathcal{I}_1 = \{\{\Phi\}, \{(L, A), (L, B)\}\}, \mathcal{I}_2 = \{\{L\}\}$$



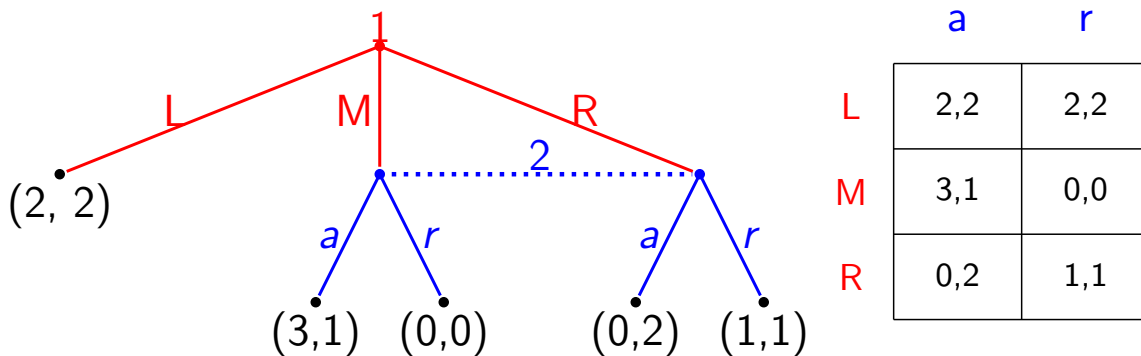
## Pure Strategies: Example 3

- ▶  $S = \{S_1, S_2\}$
- ▶  $S_1 = \{(L, a), (L, b), (R, a), (R, b)\}$
- ▶  $S_2 = \{A, B\}$



## Normal-Form Representation: Example 1

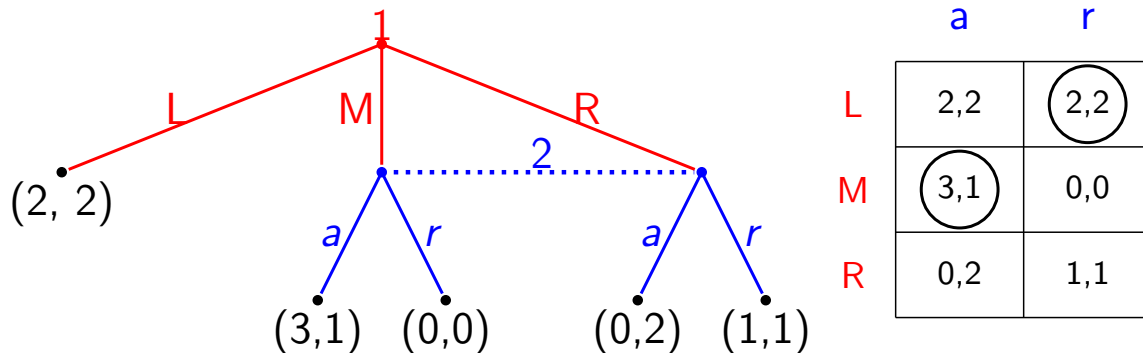
An imperfect-information extensive-form game  $\Rightarrow$  A normal-form game



The Nash Equilibrium (both pure and mixed) concept remains the same for imperfect-information extensive-form games.

## Normal-Form Representation: Example 1

An imperfect-information extensive-form game  $\Rightarrow$  A normal-form game



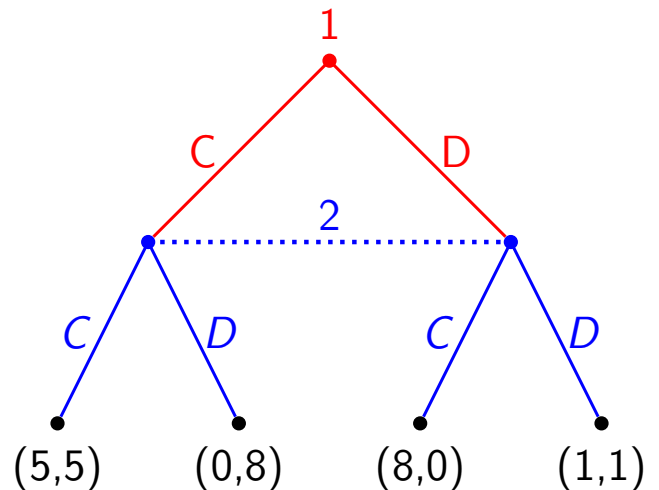
The Nash Equilibrium (both pure and mixed) concept remains the same for imperfect-information extensive-form games.

## Normal-Form Games

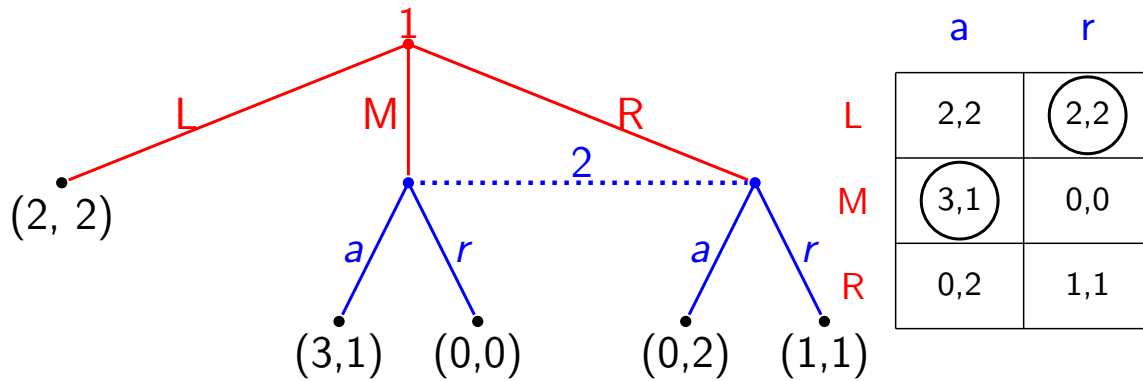
A normal-form game  $\Rightarrow$  An imperfect-information extensive-form game

	C	D
C	5, 5	0, 8
D	8, 0	1, 1

Prisoner's Dilemma

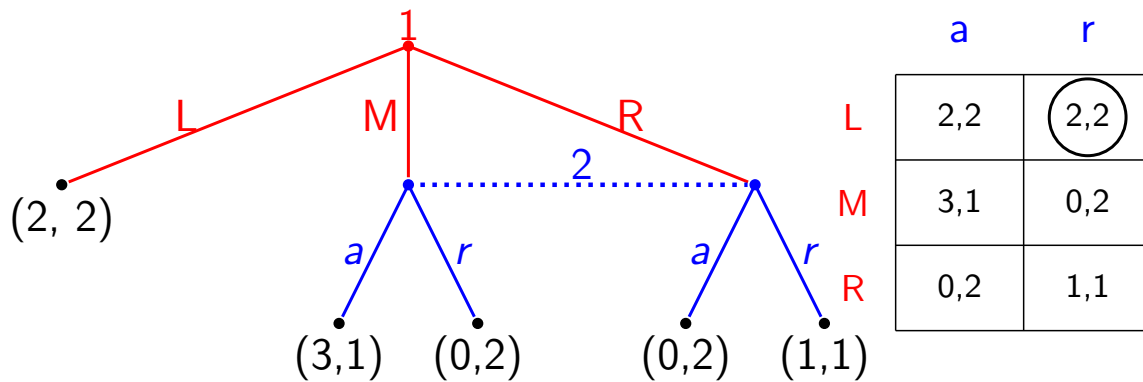


## Nash Equilibrium: Example 1



Suppose we want to generalize the idea of subgame perfect equilibrium. Consider the equilibrium (L, r). Is it subgame perfect?

## Nash Equilibrium: Example 1



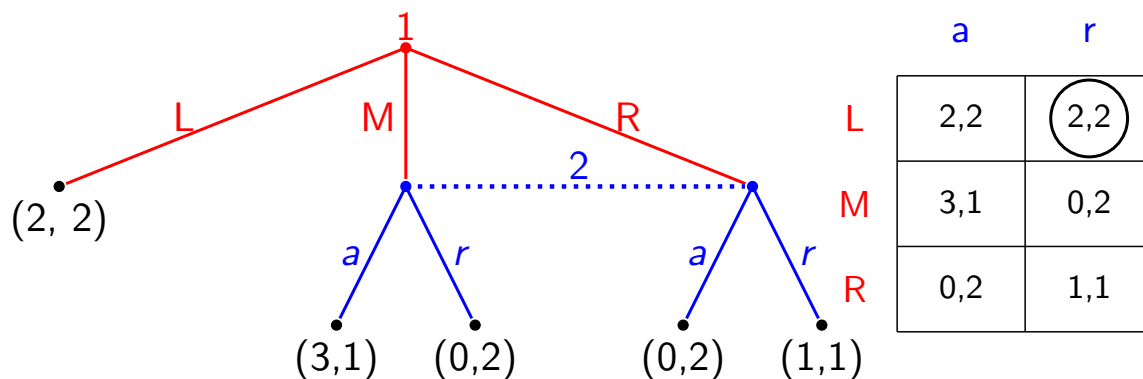
Suppose we want to generalize the idea of subgame perfect equilibrium. Consider the equilibrium (L, r). Is it subgame perfect?

## Beliefs

- ▶ A **belief**  $\mu$  is a function that assigns to every information set a probability measure on the set of histories in the information set.
- ▶ An **assessment** in an extensive-form game is a strategy-belief pair  $(s, \mu)$ .
- ▶ The assessment  $(s, \mu)$  is **sequentially rational** if for every player  $i$  and every information set  $\mathcal{I}_{i,j} \in \mathcal{I}_i$  we have

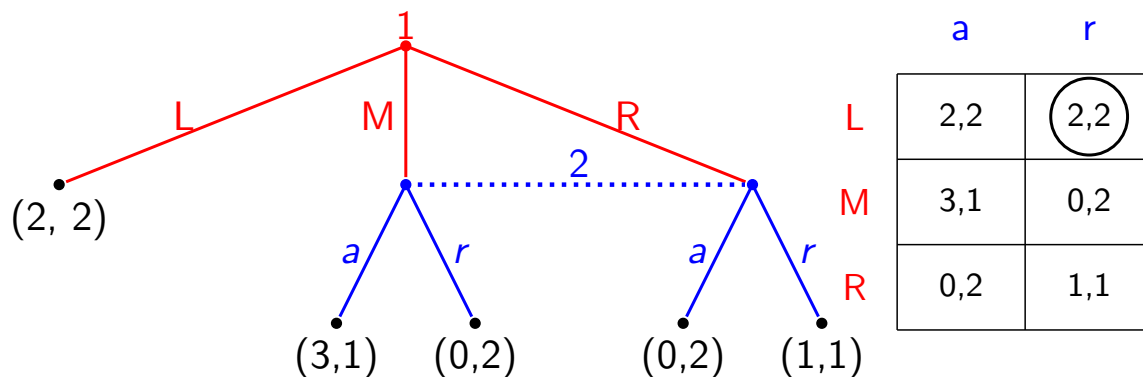
$$E[u_i(s_i, s_{-i} | \mathcal{I}_{i,j})] \geq E[u_i(s'_i, s_{-i} | \mathcal{I}_{i,j})]$$

for any  $s'_i \neq s_i$ .



## Restrictions to beliefs?

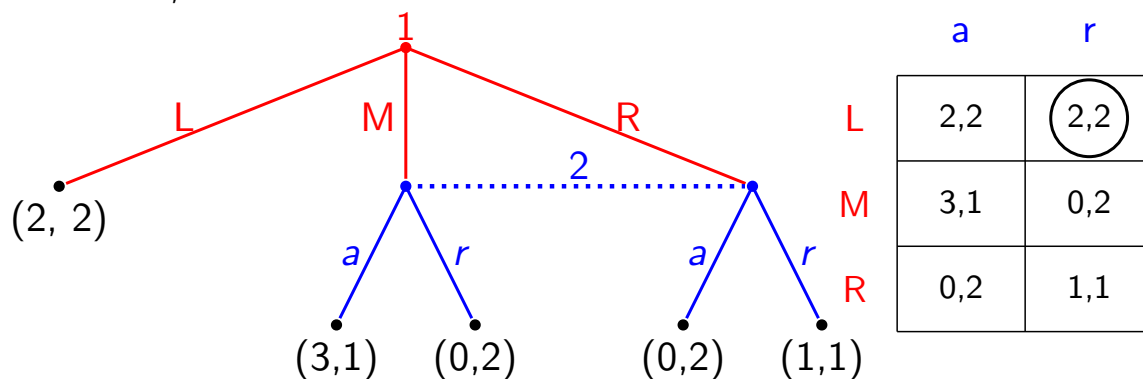
- ▶ In lieu of the Nash equilibrium concept, we require that beliefs are derived from equilibrium strategies according to Bays rule (as if players know each others strategies).
- ▶ But, what about beliefs for information sets that are **off the equilibrium path**?



## Restrictions to Beliefs?

- ▶ We want beliefs for information sets that are off the equilibrium path to be reasonable. But what is reasonable?

Consider the NE  $(L, r)$  again. Player 2's information set will not be reached at the equilibrium, because player 1 will play L with probability 1. But assume that player 1 plays a **completely mixed** strategy, playing L, M, and R with probabilities  $1 - \epsilon$ ,  $\frac{3\epsilon}{4}$ , and  $\frac{\epsilon}{4}$ . Then, the belief on player 2's information set is well defined. Now, if  $\epsilon \rightarrow 0$ , it's still well defined.



## Consistent Assessment

- ▶ An assessment  $(s, \mu)$  is **consistent** if there is a sequence  $((s^n, \mu^n))_{n=1}^{\infty}$  of assessments that converges to  $(s, \mu)$  and has the properties that each strategy profile  $s^n$  is completely mixed and that each belief system  $\mu^n$  is derived from  $s^n$  using Bayes rule.



# Sequential Equilibrium

- ▶ An assessment  $(s, \mu)$  is a **sequential equilibrium** of a finite extensive-form game with perfect recall if it is **sequentially rational** and **consistent**.
- ▶ **Thm**: Every finite extensive-form game with perfect recall has a sequential equilibrium.
- ▶ A sequential equilibrium is a Nash equilibrium.
- ▶ With perfect information, a subgame perfect equilibrium is a sequential equilibrium.

## Bayesian Games

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## So far

Up to this point, we have assumed that players know all relevant information about each other. Such games are known as games with **complete information**.

## Games with Incomplete Information

- ▶ **Bayesian Games = Games with Incomplete Information**
- ▶ **Incomplete Information**: Players have private information about something relevant to his decision making.
  - ▶ Incomplete information introduces uncertainty about the game being played.
- ▶ **Imperfect Information**: Players do not perfectly observe the actions of other players or forget their own actions.

We will see that Bayesian games can be represented as extensive-form games with imperfect information.

## Example 4: A Modified Prisoner's Dilemma Game

With probability  $\lambda$ , player 2 has the normal preferences as before (type I), while with probability  $(1 - \lambda)$ , player 2 hates to rat on his accomplice and pays a psychic penalty equal to 6 years in prison for confessing (type II).

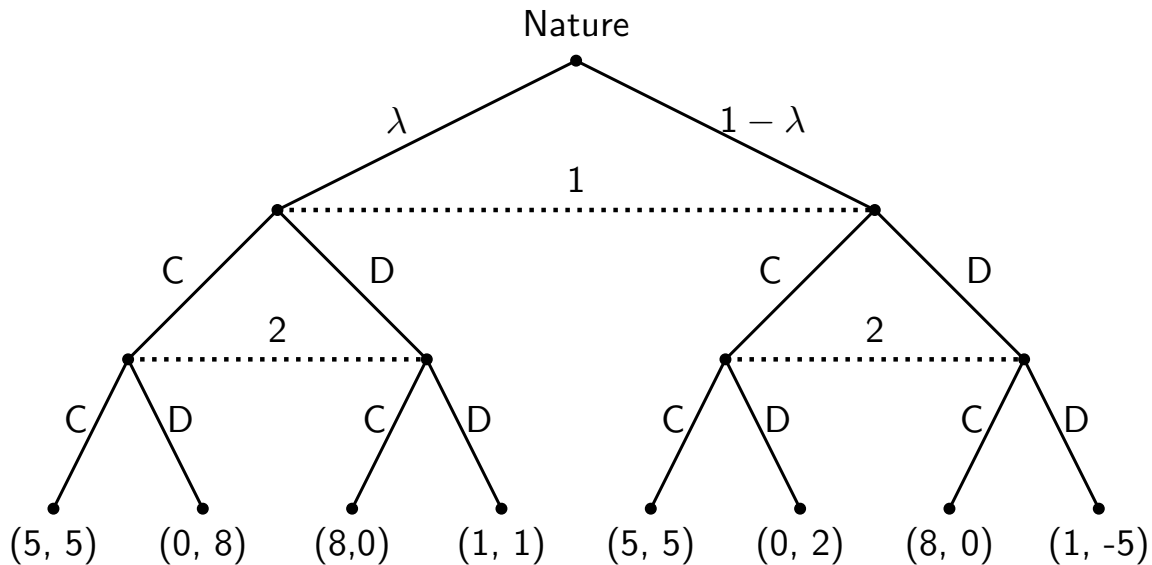
		$\lambda$	$1 - \lambda$
		C	D
C	5, 5	0, 8	5, 5
D	8, 0	1, 1	8, 0
		Type I	Type II

## Simultaneous-Move Bayesian Games

- ▶ A simultaneous-move Bayesian game is  $(N, A, \Theta, F, u)$ 
  - ▶  $N = \{1, \dots, n\}$  is the set of players
  - ▶  $A = \{A_1, A_2, \dots, A_n\}$  is the set of actions
    - $A_i = \{\text{Cooperation, Defection}\}$ .
  - ▶  $\Theta = \{\Theta_1, \Theta_2, \dots, \Theta_n\}$  is the set of types.  $\theta_i \in \Theta_i$  is a realization of types for player  $i$ .
    - $\Theta_2 = \{\text{I, II}\}$ .
  - ▶  $F : \Theta \rightarrow [0, 1]$  is a joint probability distribution, according to which types of players are drawn
    - $p(\theta_2 = \text{type I}) = \lambda$
  - ▶  $u = \{u_1, u_2, \dots, u_n\}$  where  $u_i : A \times \Theta \rightarrow \mathcal{R}$  is the utility function of player  $i$
- ▶ Two assumptions
  - ▶ All possible games have the same number of agents and the same action spaces for each agent
  - ▶ Agents have common prior. The different beliefs of agents are posteriors.

# Imperfect-Information Extensive-Form Representation of Bayesian Games

- ▶ Add a player **Nature** who has a unique strategy of randomizing in a commonly known way.



## Strategies in Bayesian Games

- ▶ A **pure strategy**  $s_i: \Theta_i \rightarrow A_i$  of player  $i$  is a mapping from every type player  $i$  could have to the action he would play if he had that type. Denote the set of pure strategies of player  $i$  as  $S_i$ .  
 $S_1 = \{\{C\}, \{D\}\}$   
 $S_2 = \{\{C \text{ if type I, C if type II}\}, \{C \text{ if type I, D if type II}\}, \{D \text{ if type I, C if type II}\}, \{D \text{ if type I, D if type II}\}\}$
- ▶ A **mixed strategy**  $\sigma_i: S_i \rightarrow [0, 1]$  of player  $i$  is a distribution over his pure strategies.

## Best Response and Bayesian Nash Equilibrium

We use pure strategies to illustrate the concepts. But they hold the same for mixed strategies.

- ▶ Player  $i$ 's **ex ante expected utility** is

$$E_{\theta}[u_i(s(\theta), \theta)] = \sum_{\theta_i \in \Theta_i} p(\theta_i) E_{\theta_{-i}}[u_i(s(\theta), \theta) | \theta_i]$$

- ▶ Player  $i$ 's **best responses** to  $s_{-i}(\theta_{-i})$  is

$$\begin{aligned} BR_i &= \arg \max_{s_i(\theta_i) \in S_i} E_{\theta}[u_i(s_i(\theta_i), s_{-i}(\theta_{-i}), \theta)] \\ &= \sum_{\theta_i \in \Theta_i} p(\theta_i) \left( \arg \max_{s_i(\theta_i) \in S_i} E_{\theta_{-i}}[u_i(s_i(\theta_i), s_{-i}(\theta_{-i}), \theta) | \theta_i] \right) \end{aligned}$$

- ▶ A strategy profile  $s_i(\theta_i)$  is a **Bayesian Nash Equilibrium** iff  $\forall i s_i(\theta_i) \in BR_i$ .

## Bayesian Nash Equilibrium: Example 4

- ▶ Playing D is a dominant strategy for type I player 2; playing C is a dominant strategy for type II player 2.
- ▶ Player 1's expected utility by playing C is  $\lambda \times 0 + (1 - \lambda) \times 5 = 5 - 5\lambda$ .
- ▶ Player 1's expected utility by playing D is  $\lambda \times 1 + (1 - \lambda) \times 8 = 8 - 7\lambda > 5 - 5\lambda$ .
- ▶ (D, (D if type I, C if type II)) is a BNE of the game.

	$\lambda$	$1 - \lambda$
	C	D
C	5, 5	0, 8
D	8, 0	1, 1
	Type I	Type II

## Example 5: An Exchange Game

- ▶ Each of two players receives a ticket  $t$  on which there is a number in  $[0,1]$ .
- ▶ The number on a player's ticket is the size of a prize that he may receive.
- ▶ The two prizes are identically and independently distributed according to a uniform distribution.
- ▶ Each player is asked independently and simultaneously whether he wants to exchange his prize for the other player's prize.
- ▶ If both players agree then the prizes are exchanged; otherwise each player receives his own prize.

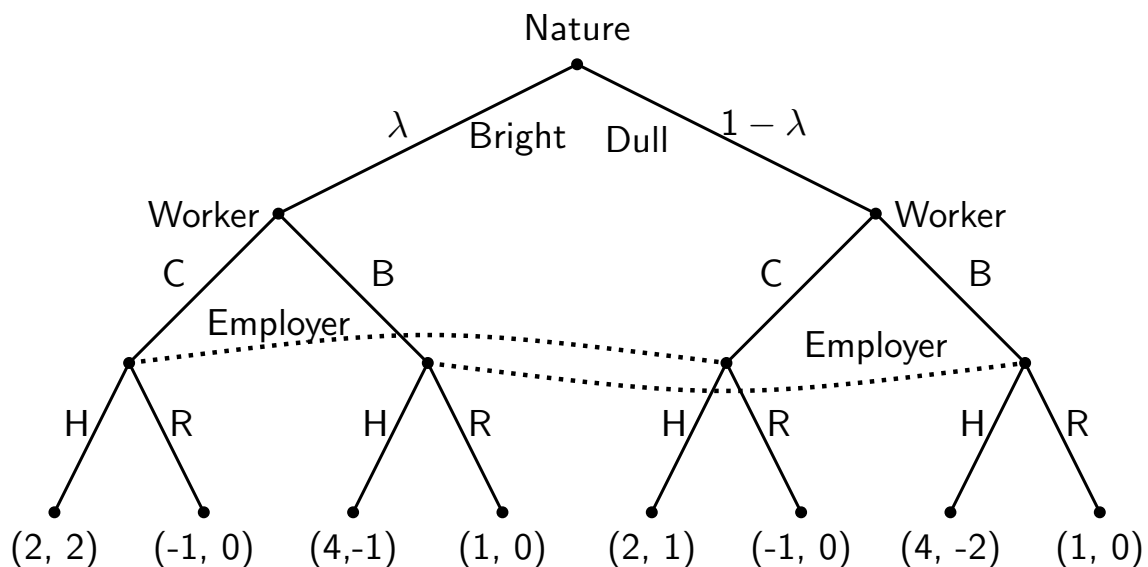
## A Bayesian Nash Equilibrium for Example 5

- ▶ Strategies of player 1 can be describe as “Exchange if  $t_1 \leq k$ ”
- ▶ Given player 1 plays such a strategy, what is the best response of player 2?
  - ▶ If  $t_2 \geq k$ , no exchange
  - ▶ If  $t_2 < k$ , exchange when  $t_2 \leq k/2$
- ▶ Since players are symmetric, player 1's best response is of the same form.
- ▶ Hence, at a Bayesian Nash equilibrium, both players are willing to exchange only when  $t_i = 0$ .

# Signaling (Sender-Receiver Games)

- ▶ There are two types of workers, bright and dull.
- ▶ Before entering the job market a worker can choose to get an education (i.e. go to college), or enjoy life (i.e. go to beach).
- ▶ The employer can observe the educational level of the worker but not his type.
- ▶ The employer can hire or reject the worker.

## Example 6: Signaling

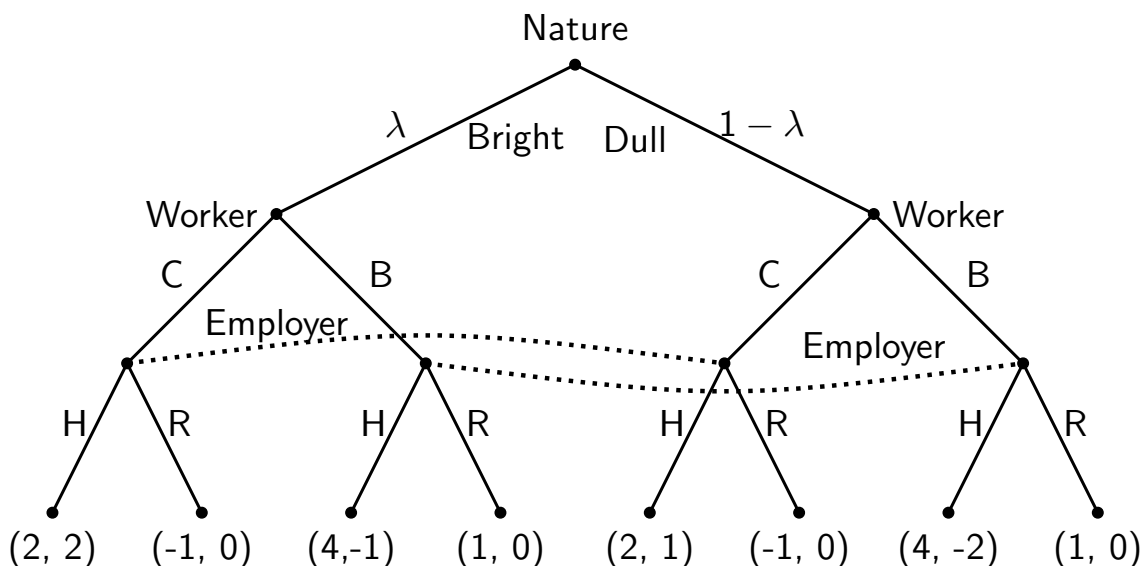


# Bayesian Extensive Games with Observable Actions

- ▶ A Bayesian extensive game with observable actions is  $(N, H, P, \Theta, p, u)$ 
  - ▶  $(N, H, P)$  is the same as those in an extensive-form game with perfect information
  - ▶  $\Theta = \{\Theta_1, \Theta_2, \dots, \Theta_n\}$  is the set of types.  $\theta_i \in \Theta_i$  is a realization of types for player  $i$ .  
 $\Theta_1 = \{\text{Bright}, \text{Dull}\}$ .
  - ▶  $F : \Theta \rightarrow [0, 1]$  is a joint probability distribution, according to which types of players are drawn  
 $p(\theta_1 = \text{Bright}) = \lambda$
  - ▶  $u = \{u_1, u_2, \dots, u_n\}$  where  $u_i : Z \times \Theta \rightarrow \mathcal{R}$  is the utility function of player  $i$ .  $Z \in H$  is the set of terminal histories.

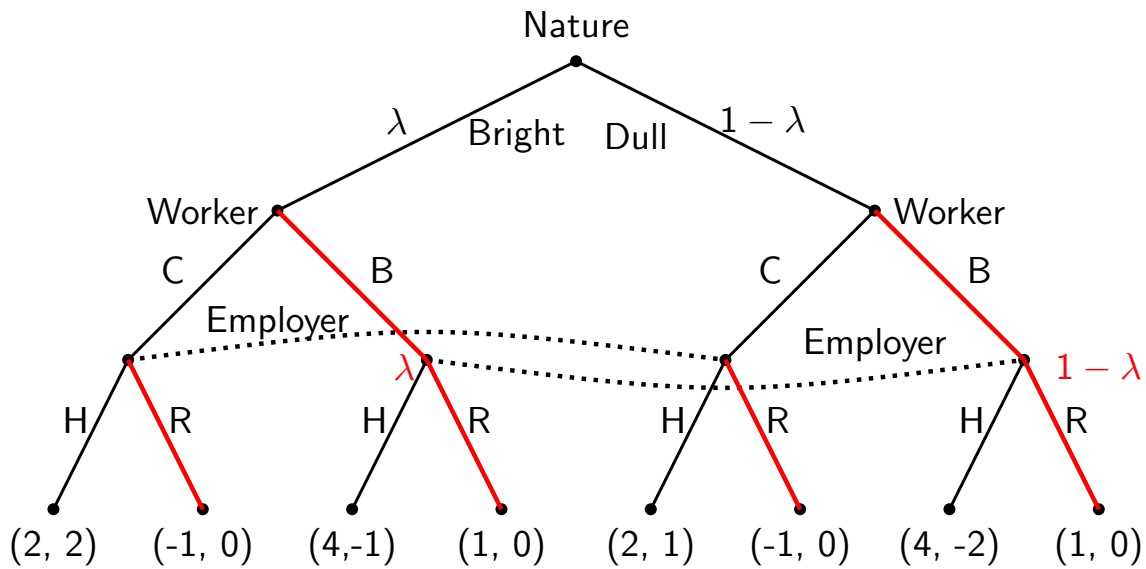
## Best Responses for Example 6

- ▶ E.g. If the employer always plays H, then the best response for the worker is B.
- ▶ But how to define best responses for the employer?
  - ▶ Beliefs on information sets
  - ▶ Beliefs derived from strategies





## A Bayesian Nash Equilibrium of Example 6



$$p(\text{Bright}|\text{Beach}) = \frac{p(\text{Bright})\sigma(\text{Beach}|\text{Bright})}{p(\text{Bright})\sigma(\text{Beach}|\text{Bright}) + p(\text{Dull})\sigma(\text{Beach}|\text{Dull})} = \frac{\lambda \cdot 1}{\lambda \cdot 1 + (1 - \lambda) \cdot 1} = \lambda$$

## “Subgame Perfection”

- ▶ The previous Bayesian Nash Equilibrium is not “subgame perfect”. When the information set College is reached, the employer should choose to hire no matter what belief he has.
- ▶ We need to require **sequential rationality** even for off-equilibrium-path information sets.
- ▶ Then, beliefs on off-equilibrium-path information sets matter.

## Perfect Bayesian Equilibrium

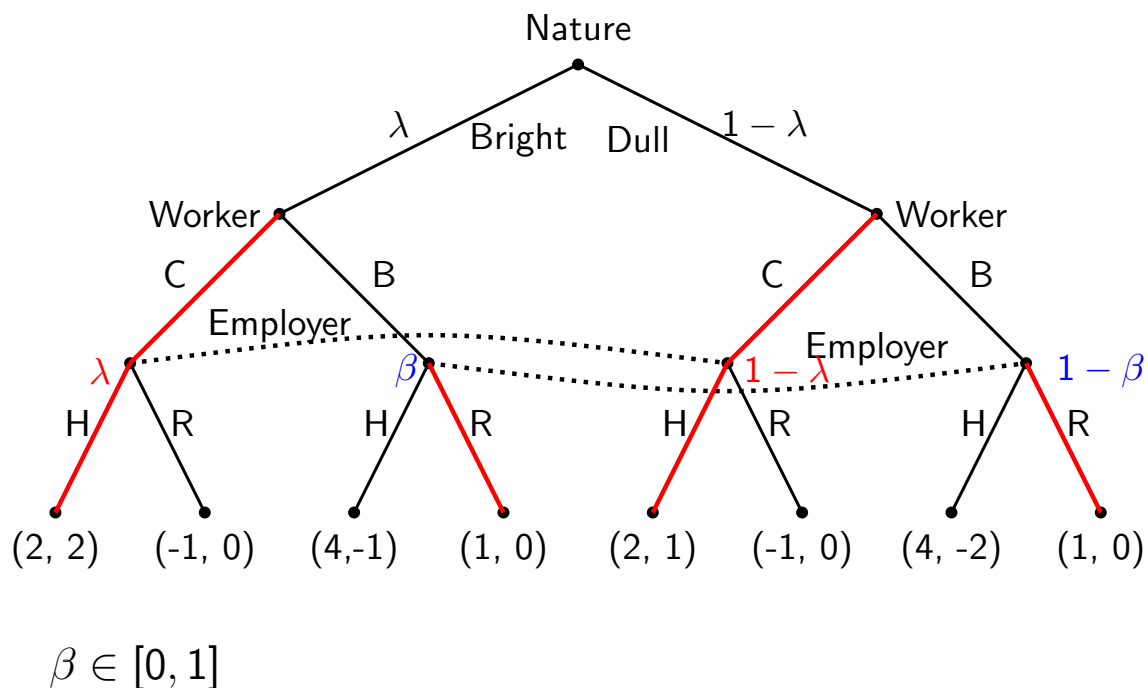
A strategy-belief pair,  $(\sigma, \mu)$  is a perfect Bayesian equilibrium if

- ▶ **(Beliefs)** At every information set of player  $i$ , the player has beliefs about the node that he is located given that the information set is reached.
- ▶ **(Sequential Rationality)** At any information set of player  $i$ , the restriction of  $(\sigma, \mu)$  to the continuation game must be a Bayesian Nash equilibrium.
- ▶ **(On-the-path beliefs)** The beliefs for any on-the-equilibrium-path information set must be derived from the strategy profile using Bayes' Rule.
- ▶ **(Off-the-path beliefs)** The beliefs at any off-the-equilibrium-path information set must be determined from the strategy profile according to Bayes Rule **whenever possible**.

## Perfect Bayesian Equilibrium

- ▶ Perfect Bayesian equilibrium is a similar concept to sequential equilibrium, both trying to achieve some sort of "subgame perfection".
- ▶ Perfect Bayesian equilibrium is defined for all extensive-form games with imperfect information, not just for Bayesian extensive games with observable actions.
- ▶ **Thm:** For Bayesian extensive games with observable actions, every sequential equilibrium is a Perfect Bayesian equilibrium.

## A Perfect Bayesian Equilibrium of Example 3



## Summary of Equilibrium Concepts

	On-equ-path strategy $\sigma_{on}$	On-equ-path belief $\mu_{on}$	Off-equ-path strategy $\sigma_{off}$	Off-equ-path belief $\mu_{off}$
NE	BR	N/A	N/A	N/A
BNE	BR given $\mu_{on}$	Consistent with $\sigma_{on}$	N/A	N/A
SPNE	BR	N/A	BR	N/A
PBE	BR given $\mu_{on}$	Consistent with $\sigma_{on}$	BR given $\mu_{off}$	Consistent with $\sigma_{off}$
SE	BR given $\mu_{on}$	Consistent with $\sigma_{on}$	BR given $\mu_{off}$	Consistent with $\sigma_{off}$