# CS161.A2

Some design considerations.

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# Design

• "Implementing to a well-designed interface."

## Design

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- You already know what your components need to do.

### Design

- "Implementing to a well-designed interface."
- You already know what your components need to do.
- Design data structures and subroutines that work towards that goal.

# File Descriptors

What happens when we open() a file?

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- The kernel chooses an index into a table of file descriptors.

• What happens when we open() a file twice?

- What happens when we open() a file twice?
- What happens when we lseek()?

- What happens when we open() a file twice?
- What happens when we lseek()?
- Two separate file descriptors

fd

• File descriptors represent open files.

#### $\mathsf{fd}$

- File descriptors represent open files.
- What represents a file?

Virtual file system

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- Abstraction layer between file-like systems (filesystem, network, devices) and OS
- structs: vnode, uio
- callables: vfs\_\*(), V0P\_\*()
- "All files are the same."

# Plan 9

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- "Everything is a file."
- Processes are files. (/proc)
- Filetypes using ioctl are just files. (/net)

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- How do you allocate new fd's (i.e., indices)?
- Should limits be set on the fd table?
  - cf. kern/include/array.h
- What do you find at fdtable[i]?

What happens when you dup2()...

- What happens when you dup2()...
  - ...and then lseek()?

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  - ...and then lseek()?
- What happens when you fork()...

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- What happens when you dup2()...
  - ...and then lseek()?
- What happens when you fork()...
  - ...and then lseek()?
  - ...and then open()?

# close()

 A file descriptor may be multiply referenced

#### close()

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- What happens if we close() a dup2()'d file descriptor?

#### close()

- A file descriptor may be multiply referenced
- What happens if we close() a dup2()'d file descriptor?
- Use pointers and reference counting.

# Processes

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  - i.e., address space + thread(s)
- Even though you're single-threading, think about the separation.
- (Nobody says you can't try multithreading...)

• What per-process state might we want?

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  - pid, address space, fd table
- What about per-thread state?
  - cwd, execution state, user stack pointer

• As with fd's, think about...

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  - ...how to allocate/free pids.
  - ...whether to enforce a limit.
  - ...modularity.

• Let's free up procs on exit().

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- What if someone is waiting on our exit status?

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- What if someone is waiting on our exit status?
- Who can wait on our exit status?

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- Who can wait on our exit status?
- Parent/child process hierarchy

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- init: ancestor of all processes

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  - ...wait for an exit()
  - ...implement WN0HANG
  - ...prevent deadlock
  - ...maintain process hierarchy w/o race

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  - address space, fd table, pid

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  - address space, fd table, pid
  - lol j/k—allocate a new pid

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- The trapframe contains execution state
  - e.g., PC and return value
- Set up process hierarchy before returning

 Replace a process's address space and execution state with a new binary

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- runprogram()

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What does runprogram() do?

- What does runprogram() do?
  - Open a binary file

- What does runprogram() do?
  - Open a binary file
  - Create a new address space

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  - Define the stack

- What does runprogram() do?
  - Open a binary file
  - Create a new address space
  - Load the executable there
  - Define the stack
  - Enter usermode

What doesn't runprogram() do?

- What doesn't runprogram() do?
  - argument handling

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  - argument handling
  - handling arguments
  - putting argv where argv has to go
  - copying arguments into your new addrspace

## Error handling

- Error handling is crucial and often hard
- What if you free your old address space, but you fail later in execv()?
- What if you fork a thread but the new thread dies before the fork finishes?

# Scheduling

#### Scheduler

- Your scheduler should be swappable with round-robin.
  - Use #ifdef's to achieve this.
- Your scheduler needs to take multiprocessing into account!

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- Timeslices increase exponentially against priority.
- Demote a thread if it uses its entire timeslice; promote it otherwise.

## Random

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- Fair, simple, low overhead, no starvation

## Lottery

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- Assign a variable number of tickets to each thread.
- Choose a random ticket and run its holder.
- Flexible, no starvation (if min 1 ticket)

#### Scheduler

- Do something better than round-robin.
- Convince us about why it's good.